SA305 - Linear Programming

# Lesson 13. Improving Search: Finding Better Solutions

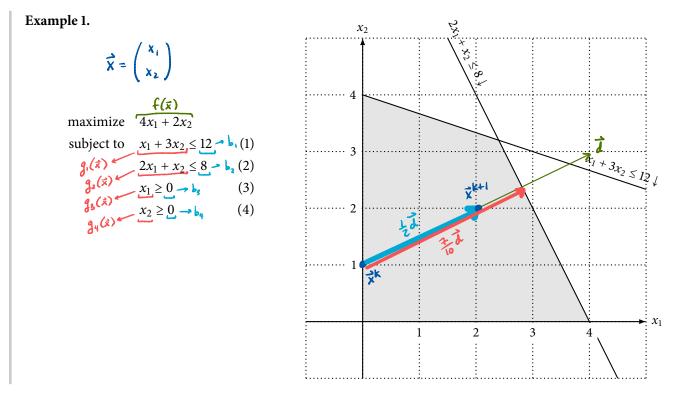
#### 1 A general optimization model

- For now, we will consider a general optimization model
- Decision variables:  $x_1, \ldots, x_n$ 
  - Recall: a feasible solution to an optimization model is a choice of values for <u>all</u> decision variables that satisfies all constraints
- Easier to refer to a feasible solution as a vector:  $\mathbf{x} = (x_1, \dots, x_n)$
- Let  $f(\mathbf{x})$  and  $g_i(\mathbf{x})$  for  $i \in \{1, ..., m\}$  be multivariable functions in  $\mathbf{x}$ , not necessarily linear
- Let  $b_i$  for  $i \in \{1, ..., m\}$  be constant scalars

minimize/maximize 
$$f(\mathbf{x})$$
  
subject to  $g_i(\mathbf{x}) \begin{cases} \leq \\ \geq \\ = \end{cases} b_i \text{ for } i \in \{1, \dots, m\}$ 

$$(*)$$

• Linear programs fit into this framework



# 2 Improving search algorithms, informally

- Idea:
  - Start at a feasible solution
  - Repeatedly move to a "close" feasible solution with better objective function value
- The neighborhood of a feasible solution is the set of all feasible solutions "close" to it
  - We can define "close" in various ways to design different types of algorithms
- Let's start formalizing these ideas

#### 3 Locally and globally optimal solutions

•  $\varepsilon$ -neighborhood  $N_{\varepsilon}(\mathbf{x})$  of a solution  $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{R}^n$  (where  $\varepsilon > 0$ ):

$$N_{\varepsilon}(\mathbf{x}) = \{\mathbf{y} \in \mathbb{R}^n : d(\mathbf{x}, \mathbf{y}) \leq \varepsilon\}$$

where  $d(\mathbf{x}, \mathbf{y})$  is the distance between solution  $\mathbf{x}$  and  $\mathbf{y}$ 

• A feasible solution **x** to optimization model (\*) is **locally optimal** if for some value of  $\varepsilon > 0$ :

 $f(\mathbf{x})$  is better than  $f(\mathbf{y})$  for all feasible solutions  $\mathbf{y} \in N_{\varepsilon}(\mathbf{x})$ 

• A feasible solution **x** to optimization model (\*) is **globally optimal** if:

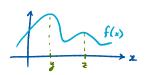
 $f(\mathbf{x})$  is better than  $f(\mathbf{y})$  for all feasible solutions  $\mathbf{y}$ 

• Also known simply as an **optimal solution** 

- Global optimal solutions are locally optimal, but not vice versa
- In general: harder to check for global optimality, easier to check for local optimality

### 4 The improving search algorithm

- 1: Find an initial feasible solution  $\mathbf{x}^0$
- 2: Set k = 0
- 3: while  $\mathbf{x}^k$  is <u>not</u> locally optimal **do**
- 4: Determine a new feasible solution  $\mathbf{x}^{k+1}$  that improves the objective value at  $\mathbf{x}^k$
- 5: Set k = k + 1
- 6: end while
- Generates sequence of feasible solutions  $\mathbf{x}^0, \mathbf{x}^1, \mathbf{x}^2, \dots$
- In general, improving search converges to a local optimal solution, not a global optimal solution
- Let's concentrate on line 4 finding better feasible solutions



 $\vec{x}^{k+1} = \vec{x}^k + \lambda \vec{d} + feasible$ 

### 5 Moving between solutions

• How do we move from one solution to the next?

to the next?  
new 
$$\mathbf{x}^{k+1} = \mathbf{x}^k + \lambda \mathbf{d}$$
 direction  
solution  $\mathbf{x}^{k+1} = \mathbf{x}^k + \lambda \mathbf{d}$ 

solution

• In Example 1:

Let 
$$\vec{x}^{k} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$
  $\lambda = \frac{1}{2}$   $\vec{d} = \begin{pmatrix} 4 \\ 2 \end{pmatrix}$   
 $\Rightarrow \vec{x}^{k+1} = \vec{x}^{k} + \lambda \vec{d} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} + \frac{1}{2} \begin{pmatrix} 4 \\ 2 \end{pmatrix} = \begin{pmatrix} 2 \\ 2 \end{pmatrix}$ 

# 6 Improving directions

- We want to choose **d** so that  $\mathbf{x}^{k+1}$  has a better value than  $\mathbf{x}^k$
- **d** is an **improving direction** at solution  $\mathbf{x}^k$  if

$$f(\mathbf{x}^k + \lambda \mathbf{d})$$
 is better than  $f(\mathbf{x}^k)$ 

$$\vec{a} = \begin{pmatrix} a_1 \\ a_2 \end{pmatrix} \qquad \vec{b} = \begin{pmatrix} b_1 \\ b_2 \end{pmatrix}$$
$$\vec{a} \cdot \vec{b} = a_1 b_1 + a_2 b_2$$
$$\vec{a}^{\mathsf{T}} \vec{b} = a_1 b_1 + a_2 b_2$$

for all positive 
$$\lambda$$
 "close" to 0

- How do we find an improving direction?
- The **directional derivative** of f in the direction **d** at solution  $\mathbf{x}^k$  is

$$\frac{\nabla f(\vec{x}^k)^T \vec{d}}{\|\vec{d}\|} = \text{rate of change in } f \text{ at } \vec{x}^k \text{ in the direction } \vec{d}$$

• Maximizing  $f: \mathbf{d}$  is an improving direction at  $\mathbf{x}^k$  if

$$\nabla f(x^{*})^{T} \vec{a} > 0$$

• Minimizing  $f: \mathbf{d}$  is an improving direction at  $\mathbf{x}^k$  if

 $\nabla f(\vec{x}^{k})^{T}\vec{a} < 0$ 

• In Example 1:

Is 
$$\vec{d} = \begin{pmatrix} 4 \\ 2 \end{pmatrix}$$
 improving at  $\vec{x}^{k} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$ ?  
 $f(\vec{x}) = 4x_{1} + 2x_{2} \implies \nabla f(\vec{x}) = \begin{pmatrix} 4 \\ 2 \end{pmatrix}$  (for any  $\vec{x}^{!}$ )  
 $\Rightarrow \nabla f(\vec{x}^{k})^{T} \vec{d} = \begin{pmatrix} 4 \\ 2 \end{pmatrix}^{T} \begin{pmatrix} 4 \\ 2 \end{pmatrix} = 20 > 0$   
 $\Rightarrow Y_{40}, \vec{d}$  is improving at  $\vec{x}^{k}$   
 $= 4(4) + 2(2)$   
 $= 20$ 

• For linear programs in general: if **d** is an improving direction at  $\mathbf{x}^k$ , then  $f(\mathbf{x}^k + \lambda \mathbf{d})$  improves as  $\lambda \to \infty$ 

# 7 Step size

- We have an improving direction **d** now how far do we go?
- One idea: find maximum value of  $\lambda$  so that  $\mathbf{x}^k + \lambda \mathbf{d}$  is still feasible
- Graphically, we can eyeball this
- Algebraically, we can compute this in Example 1:

For what values of 
$$\lambda$$
 is  $\vec{x}^{\pm} + \lambda \vec{\lambda} = \begin{pmatrix} 4\lambda \\ 1+2\lambda \end{pmatrix}$  featible?  

$$\frac{f(\vec{x})}{4x_1+2x_2}$$
subject to  $x_1 + 3x_2 \leq 12^{-1} \cdot (.0)$ 

$$g(\vec{x}) = \frac{2x_1 + x_2}{2x_2 + x_2} \leq 8^{-1} \cdot (.2)$$

$$g(\vec{x}) = \frac{2x_1 + x_2}{2x_2 + x_2} \leq 8^{-1} \cdot (.2)$$

$$g(\vec{x}) = \frac{2x_1 + x_2}{2x_2 - 1} \leq (.2)$$

$$(1) \quad x_1 + 3x_2 \leq 12$$

$$(2) \quad 2x_1 + x_2 \leq 8$$

$$(3) \quad 4\lambda + 3(1+2\lambda) \leq 12$$

$$(2) \quad 2x_1 + x_2 \leq 8$$

$$(3) \quad 4\lambda + 3(1+2\lambda) \leq 12$$

$$(4) \quad x_1 \leq 3$$

$$(5) \quad x_1 \geq 0$$

$$(4) \quad x_1 \geq 0$$

$$(5) \quad x_1 \geq 12$$

$$(5) \quad x_$$

#### 8 Feasible directions

- Some improving directions don't lead to any new feasible solutions
- **d** is a **feasible direction** at feasible solution  $\mathbf{x}^k$  if  $\mathbf{x}^k + \lambda \mathbf{d}$  is feasible for all positive  $\lambda$  "close" to 0
- Again, graphically, we can eyeball this
- A constraint is **active** at feasible solution **x** if it is satisfied with equality

• For linear programs:

• We have constraints of the form:

$$a_1x_1 + a_2x_2 + \dots + a_nx_n \le b$$
$$a_1x_1 + a_2x_2 + \dots + a_nx_n \ge b$$
$$a_1x_1 + a_2x_2 + \dots + a_nx_n = b$$

• We can rewrite these constraints using vector notation:

Let: 
$$\vec{a} = \begin{pmatrix} a_1 \\ a_2 \\ \vdots \\ a_n \end{pmatrix}$$
  $\vec{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}$   $\Rightarrow$   $\vec{a}^T \vec{x} \ge b$   
 $\vec{a}^T \vec{x} = b$ 

• **d** is a feasible direction at **x** if

 $\diamond \ \mathbf{a}^{\mathsf{T}}\mathbf{d} \leq 0 \text{ for each <u>active</u> constraint of the form } \mathbf{a}^{\mathsf{T}}\mathbf{x} \leq b$ 

 $\diamond \mathbf{a}^{\mathsf{T}}\mathbf{d} \ge 0$  for each active constraint of the form  $\mathbf{a}^{\mathsf{T}}\mathbf{x} \ge b$ 

 $\mathbf{a}^{\mathsf{T}}\mathbf{d} = 0$  for each active constraint of the form  $\mathbf{a}^{\mathsf{T}}\mathbf{x} = b$ 

If there are no active constraints at a feasible solution x, then any direction is feasible.

• In Example 1:

Is  $\vec{d} = \begin{pmatrix} 4 \\ 2 \end{pmatrix}$  a feasible direction at  $\vec{x}^{k} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$ ? f(x) maximize  $4x_1 + 2x_2$ subject to  $x_1 + 3x_2 \le 12$  **b** (1)  $g_{1}(\vec{x}) = \frac{2x_{1} + x_{2}}{y_{1}(\vec{x})} \leq \frac{2x_{1} + x_{2}}{x_{1} > 0} \leq \frac{8}{x_{1}} = b_{2}(2)$ Active constraints?  $\begin{array}{c}
 g_{1}(\vec{x}) & \underbrace{x_{1} \geq 0}_{g_{1}(\vec{x})} & \underbrace{x_{1} \geq 0}_{g_{1}(\vec{x})} & \underbrace{x_{1} \geq 0}_{g_{2}(\vec{x})} & \underbrace{x_{2} \geq 0}_{g_{4}(\vec{x})} & \underbrace{x_{2} \geq 0}_{g_{4}(\vec{$ (3) (1)  $0 + 3(1) \le 12$  not active (4)not active  $(2) 2(0) + (1) \leq 8$ active → (3) 0 2 0 not active 1 2 0 (4) (3) is of the form  $\vec{a}^T \vec{x} \ge 0$   $\begin{pmatrix} 1 \\ 0 \end{pmatrix}^T \begin{pmatrix} x_1 \\ x_2 \end{pmatrix}$   $\vec{a}$  is feasible at  $\vec{x}^k$ . => need to check if and = 0 => Yes!  $\binom{1}{2}^{T}\binom{4}{2}$ 

# 9 Detecting unboundedness - improve obj. fr. value infinitely.

- Suppose **d** is an improving direction at feasible solution  $\mathbf{x}^k$  to a <u>linear</u> program
- Also, suppose  $\mathbf{x}^k + \lambda \mathbf{d}$  is feasible for all  $\lambda \ge 0$
- What can you conclude?

| LP is unbounded : | $f(\vec{x}^{k}+\lambda \vec{d})$ improves and $\vec{x}^{k}+\lambda \vec{d}$ |
|-------------------|---|
|                   | remains feasible as $\lambda \rightarrow \infty$                            |
|                   |   |
|                   |   |

# 10 Summary

- Line 4 boils down to finding an improving and feasible direction **d** and an accompanying step size  $\lambda$
- We discussed conditions on whether a direction is improving and feasible
- We don't know how to systematically find such directions... yet